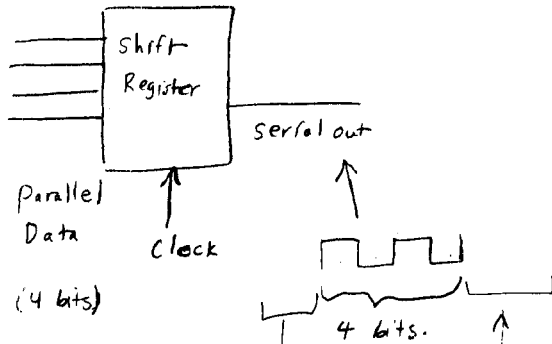


# Sending PCM

Data bits are set in serial.  
 Parity bits might be added  
 Sync bits must be added



Add a "sync" bit ("Framing" bit in text)  
 This is something easily recognizable —  
 Always 1?

For better reliability —  
 use a known sequence  
 not just a single bit.

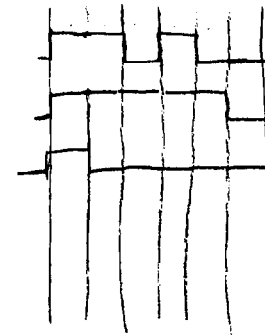
Add a "stop bit" —  
 really just what happens when data  
 runs out — but you need to  
 account for it.  
 Needs to be easily recognizable —  
 always 0?

Hw: 6.2-1, 2, 3, 4, 6, 7

10B-1

So --- to send  
 1010  
 1111  
 0000

Waveform is:



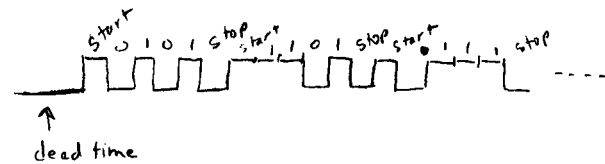
10B-2

This system has 2 bits overhead for every sample —  
 "binary code word"

So --- to send:

5, 13, 7, 4, 0:

0101, 1101, 0111, 0100, 0000



For efficiency — use long word lengths:

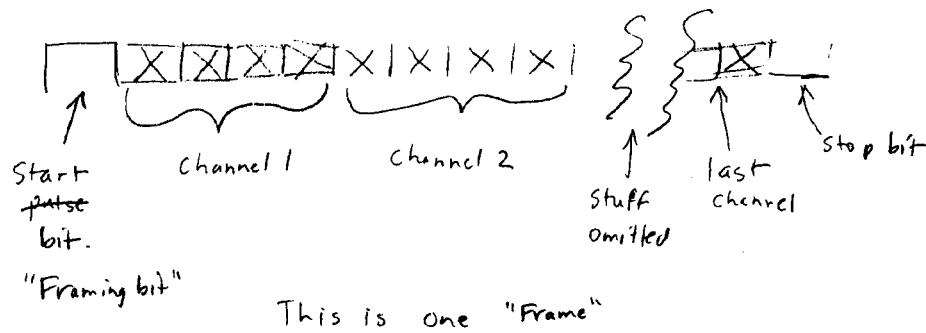
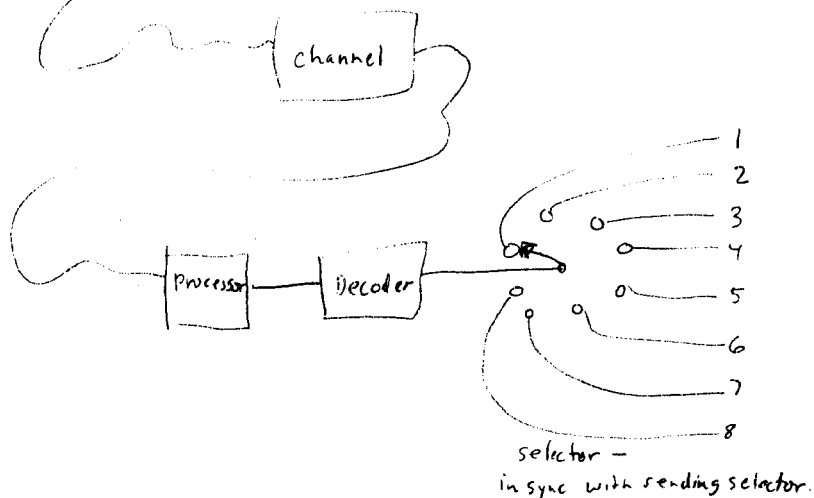
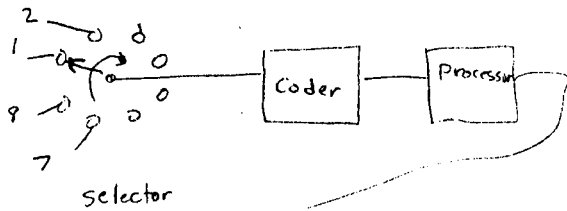
so start & stop take less time, by percentage.

or --- send multiple code-words with single sync.

# Time division multiplexing -

10B-3

intermix the codes of multiple signals.



# Bandwidth ---

10B-4

if our channel has bandwidth  $B$  - A cycle can transmit  $2B$  bits  
so. can transmit  $2B$  samples per second.

To send a signal with  $n$  bits of bandwidth  $B$  -  
need  $nB$  bits per second.

This "n bits" includes framing bits.

Example: - We want to send 16 3.5 kHz  
voice channels (telephone quality)

with 30 dB S/N ratio, using PCM, TDM.  
assuming square wave modulation, sample 20% above  
Nyquist rate.

Find # of samples, bandwidth, etc.

For 30 dB S/N --

$$\frac{S}{N} = 3L^2 \frac{\overline{m^2(t)}}{m_p^2}$$

For square wave --

$$\frac{\overline{m^2(t)}}{m_p^2} = 1$$

$$\frac{S}{N}_{dB} = 30 \text{ dB} \Rightarrow \frac{S}{N}_{\text{Power}} = 1000$$

$$1000 = 3L^2$$

$$L^2 = \frac{1000}{3}$$

$$L = 18.2574$$

Need 19 levels.

For power of 2, use 32 levels, 5 bits  
will result in  $S_r = 3072$  or 24.97 dB

check:  $10 \log_{10}(3) + 6n = 34.97 \text{ dB}$   
could have solved this way.

For  $B_{\text{band}} = 3.5 \text{ kHz}$

10B-5

Nyquist rate is 7 kHz.

20% above is: 8.4 kHz.

For 16 channels, 5 bits each -

need  $16 \times 5 = 80$  bits.

+ 2 bits for framing

= 82 bits per frame.

8.4 kHz sample rate  $\Rightarrow 1.19 \times 10^{-4} \frac{\text{sec}}{\text{sample}}$ .

We need to send 82 bits in  $1.19 \times 10^{-4} \text{ sec}$ .

$1.452 \times 10^{-6} \text{ sec/bit}$

Data clock = 688.8 kHz.

We can send 2 bits per cycle -

so channel bandwidth = 344.4 kHz.

One channel would be -

Need 7 bits in  $1.19 \times 10^{-4} \text{ sec}$

=  $1.7 \times 10^{-5} \text{ sec/bit}$

Data clock -  $5.882 \times 10^4 = 58.8 \text{ kHz}$

channel BW = 29.41 kHz

Text: sec. 6.2.4 describes a T1 (telephone) system.

$\frac{S}{N}$  ratio with nonuniform quantization 10B-6

Without deriving -

$$\frac{S_0}{N_0} = C 2^{2n} = C L^2$$

$n = \# \text{ of bits}$   
 $2^n = L = \# \text{ of levels}$   
 $2^{2n} = L^2$

$$C = \frac{3}{[\ln(1+u)]^2} \leftarrow \mu \text{ law encoding}$$

when  $\frac{m_p^2}{m^2(f)} \ll u^2$

5 bits,  $\mu = 10$  would give us:

$$\frac{3}{[\ln(11)]^2} = 0.52$$

$$\frac{S}{N} = 0.52 (32)^2 = 532.4 = 27.26 \text{ dB}$$

